

FIG. 1

K=3, l=2 convolutional encoder

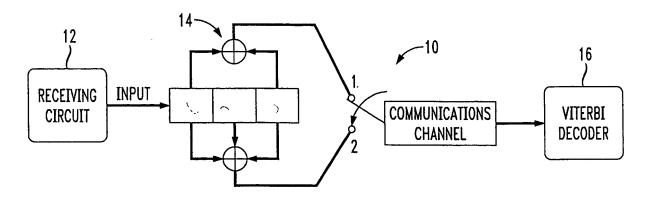
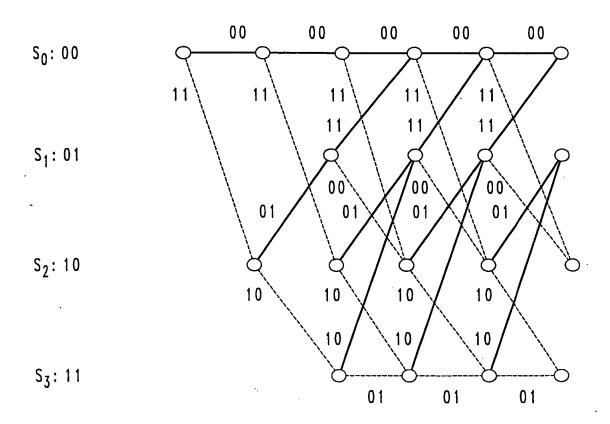


FIG. 2

Trellis for ordinary convolutional code

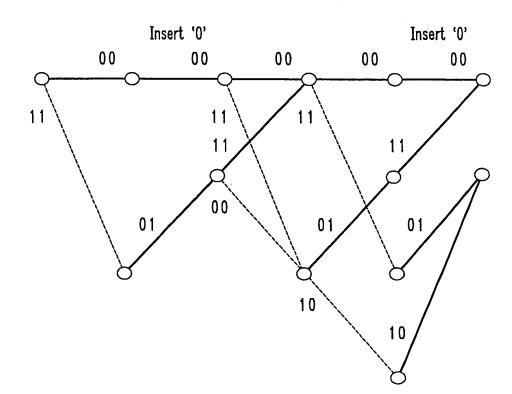


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FIG. 3

Inserting zero at the first position periodically



S₀: 00

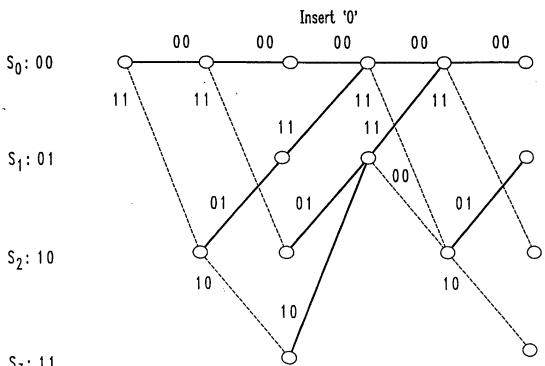
S₁: 01

S₂: 10

S₃: 11

FIG. 4

Insert zero at the second position periodically



S₇: 11

FIG. 5

Generator Matrix

Let

$$C(j) = X(j)G, \quad j = 1, 2, ..., K-1,$$
 (1)

where $X^{(j)} = [1, x_1, ..., x_{j-1}, 0, x_{j+1}, ...], x_{tK+j} = 0, t = 0, 1, ..., G$ is the Toeplitz block matrix

$$G = [\vec{g}_{i-j}]_{i,j=0,1,...}$$

with 1 \times K sub-blocks

$$\vec{g}_{i} = \begin{cases} [g_{1,i}, g_{2,i}, ..., g_{l,i}], & i = 0,1, ..., m; \\ 0, & \text{others.} \end{cases}$$



FIG. 6

Gj Presentation

$$\begin{bmatrix} \overrightarrow{g}_0(t) & \overrightarrow{g}_1(t+1) & \cdots & \overrightarrow{g}_m(t+m) & \cdots \\ 0 & \overrightarrow{g}_0(t+1) & \cdots & \overrightarrow{g}_{m-1}(t+m) & \cdots \\ \vdots & \vdots & \ddots & \vdots \end{bmatrix},$$

FIG. 6A

$$\begin{aligned} \phi_t \left(X_{t-K+2}^t \right) \\ &= \max_{\substack{X \ t-K+1 \\ 0}} M \left(X_0^t \right) \\ &= \max_{\substack{x \ t-K+1}} \left[L \left(X_{t-K+1}^t \right) + \phi_{t-1} \left(X_{t-K+1}^{t-1} \right) \right] \end{aligned}$$

FIG. 7A

DECODING

Step 1 Initialization: For $0 \le t < K-1$, start—ing from $\phi\left(X_{-K}^{-1}\right) = 0$ we calculate $\phi\left(X_{t-K+1}^{t}\right)$ for all possible combinations of X_{0}^{t} by (3).

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Step 2 Recursive forward algorithm at t:

If $t \neq K-1 \pmod K$, we compute $\phi\left(X_{t-K+2}^t\right)$ by (3) and save

$$\widetilde{x}_{t-K+1} \left(X_{t-K+2}^t \right)$$

$$= \underset{x_{t-K+1}}{\arg \max} \left[L \left(X_{t-K+1}^t \right) + \phi \left(X_{t-K+1}^{t-1} \right) \right]; (5)$$

otherwise we compute $\phi\left(X_{t-K+2}^{t}\right)$ by (4). Go to Step 3.

FIG. 7B

Step 3 Recursive backward algorithm at t

If $t - D \neq K - 1 \pmod{K}$, starting from

$$\hat{X}_{t-K+2}^{t} = \underset{X_{t-K+2}}{\text{arg max }} \phi \left(X_{t-K+2}^{t} \right)$$
 (6)

we calculate $\hat{x}_k = \tilde{x}_k \left(\hat{X}_{k+1}^{k+K-1} \right)$, k = t-K+1, t-

 $K, t-K-1, \ldots$ until backward D symbols to find

$$\hat{x}_{t-D} = \tilde{x}_{t-D} \left(\hat{X}_{t-D+1}^{t-D+K-1} \right);$$
 (7)

otherwise $\hat{x}_{t-D} = 0$.

T ≠ N, Back 34 to Step 2

If t = n go to Step 4; otherwise go to Step 2.

Step 4 Termination: Let $n \le t < n + K - 2 = N$.

If $t \neq K-1 \pmod K$, we compute $\phi\left(X_{t-K+2}^t\right)$ by (3) and save $\widetilde{x}_{t-K+1}\left(X_{t-K+2}^t\right)$ by (5); otherwise

we compute $\phi\left(X_{t-K+2}^{t}\right)$ by (4) and we do not need

to save \tilde{x}_{t-K+1} $\left(X_{t-K+2}^{t} \right)$ since it must be zero.

Repeat this step until t = N, then go to Step 5.

FIG. 7C

From Step 4

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Step 5 Recursive backward algorithm at the end: Starting from

$$\hat{x}_n = \arg \max_{x_n} \phi \left(\underbrace{0, ..., 0}_{K-2}, x_n \right),$$

we estimate $\boldsymbol{x_t}$ by

$$\hat{x}_t = \tilde{x}_t \left(\hat{X}_t^{t+K-2} \right), \ t = n-1, n-2, ..., n-D.$$

FIG. 8

Code	Conv. Code	Conv. Zero Code
Code Rate	$\frac{T}{(T+K-1)l}\approx\frac{1}{l}$	$\frac{T}{Nl} \approx \frac{K-1}{Kl}$
Complexity	$\approx T(l+2)2^K$	$\approx \frac{K}{K-1} T(l+2) 2^{K-1}$
Memory	$_{2}K_{D}$	$2^{K-1}\left(D-\left[\frac{D}{K}\right]\right)$
Delay	D	D